Lesson 12. Dynamic Programming - Review

- Recall from Lessons 5-11:
 - A **dynamic program** models situations where decisions are made in a <u>sequential</u> process in order to optimize some objective
 - **Stages** t = 1, 2, ..., T
 - \diamond stage *T* \leftrightarrow end of decision process
 - **States** $n = 0, 1, ..., N \leftarrow$ possible conditions of the system at each stage
 - Two representations: **shortest/longest path** and **recursive**

Shortest/longest path		Recursive	
node t_n	\leftrightarrow	state <i>n</i> at stage <i>t</i>	
$edge(t_n,(t+1)_m)$	\leftrightarrow	allowable decision x_t in state n at stage t that results in being in state m at stage $t + 1$	
length of edge $(t_n, (t+1)_m)$	\leftrightarrow	cost/reward of decision x_t in state n at stage t that results in being in state m at stage $t + 1$	
length of shortest/longest path from node t_n to end node	\leftrightarrow	cost/reward-to-go function $f_t(n)$	
length of edges (T_n, end)	\leftrightarrow	boundary conditions $f_T(n)$	
shortest or longest path	\leftrightarrow	recursion is min or max:	
		$f_t(n) = \min_{x_t \text{ allowable}} \left\{ \begin{pmatrix} \operatorname{cost/reward of} \\ \operatorname{decision} x_t \end{pmatrix} + f_{t+1} \begin{pmatrix} \operatorname{new state} \\ \operatorname{resulting} \\ \operatorname{from} x_t \end{pmatrix} \right\}$	
source node 1 _n	\leftrightarrow	desired cost-to-go function value $f_1(n)$	

Oil capacity per day	Gas capacity per day	Building cost (\$ millions)
0	0	0
1000	0	5
0	1000	7
1000	1000	14

Example 1. Simplexville Oil needs to build capacity to refine 1,000 barrels of oil and 2,000 barrels of gasoline per day. Simplexville can build a refinery at 2 locations. The cost of building a refinery is as follows:

The problem is to determine how much capacity should be built at each location in order to minimize the total building cost. To make things a little simpler, assume that the capacity requirements must be met exactly.

- a. Formulate this problem as a dynamic program by giving its shortest path representation.
- b. Formulate this problem as a dynamic program by giving its recursive representation. Solve the dynamic program.

Stage t represent deciding to build at location t (t=1,2) or the end of the decision-making process (t=3)

State
$$(n_1, n_2)$$
 represents having n_1 oil capacity and n_2 gas capacity still needed to be built $(n_1 = 0, 1; n_2 = 0, 1, 2)$



Find the shortest path:

Recursive representation

- Stage t represents deciding to build at location t (t=1,2) or the end of the decision-making process (t=3)
- · State (n, n2) represents having n, ril capacity and n2 gas capacity still needed to be built $(n_1 = 0, 1; n_2 = 0, 1, 2)$
- · Allowable decisions xt at staget and state (n,, n2):

$$t=1,2: \quad x_{t} = (x_{t_{1}}, x_{t_{2}}) = build x_{t_{1}} \text{ oil capacity and } x_{t_{2}} \text{ gas capacity at location } t$$

$$x_{t} \quad must \quad satisfy: \quad x_{t_{1}} \in \{0,1\}$$

$$x_{t_{2}} \in \{0,1\}$$

$$x_{t_{1}} \in \{0,1\}$$

$$x_{t_{1}} \in \{0,1\}$$

$$x_{t_{1}} \in \{0,1\}$$

$$x_{t_{2}} \in \{0,1\}$$

t=3: no decisions

• Cost of
$$z_t$$
 at stage t and state (n_1, n_2) :
• Cost of z_t at stage t and state (n_1, n_2) :
• $C(x_{t1}, x_{t2}) = \begin{cases} 0 & \text{if } (x_{t1}, x_{t2}) = (0, 0) \\ 5 & \text{if } (x_{t1}, x_{t2}) = (1, 0) \\ 7 & \text{if } (x_{t1}, x_{t2}) = (0, 1) \\ 14 & \text{if } (x_{t1}, x_{t2}) = (1, 1) \end{cases}$
• for $t = 1, 2$
• $n_1 = 0, 1$
• $n_2 = 0, 1, 2$

- · Cost-to-go function: $f_t(n_1, n_2) = minimum cost to build n_1 oil and n_2 gas capacity$ for t = 1, 2, 3 $n_{1} = 0, 1$ with locations t, t+1, ... available $n_2 = 0, 1, 2$
- Boundary conditions: $f_3(n_1, n_2) = \begin{cases} 0 & \text{if } (n_1, n_2) = (0, 0) \\ +\infty & 0 \\ \end{vmatrix}$ for N, = 0,1 $N_2 = 0, 1, 2$
- Remission: $\begin{aligned}
 & f_t(n_1, n_2) = \min_{\substack{x_{t1} \in \{0, 1\}\\x_{t2} \in \{0, 1\}}} \begin{cases} C(x_{t1}, x_{t2}) + f_{t+1}(n_1 x_{t1}, n_2 x_{t2}) \\ n_1 = 0, 1 \\ n_2 = 0, 1, 2 \end{cases} \quad f_{t1} = 0, 1, 2
 \end{aligned}$ $\chi_{t1} \leq n_1, \chi_{t2} \leq n_2$
- · Desired cost-to-go value: f. (1,2)

Stage 3:
(boundary)
$$f_3(n_1, n_2) = \begin{cases} 0 & \text{if } (n_1, n_2) = (0, 0) \\ +\infty & 0 \\ \end{pmatrix} \quad for \quad n_1 = 0, 1 \\ n_2 = 0, 1, 2 \\ n_2 = 0, 1, 2 \end{cases}$$

Stage 2:

$$f_{2}(1,2) = \min \begin{cases} 0+\infty = +\infty & f_{2}+\infty = +\infty & f_{1}+\infty = +\infty \\ f_{2}(1,2) = \min \begin{cases} 0+\infty = +\infty & f_{3}(1,2), c(1,0) + f_{3}(0,2), c(0,1) + f_{3}(1,1), c(1,1) + f_{3}(0,1) \\ (0,0) & (0,1) & (0,1) \\ f_{2}(1,1) = \min \begin{cases} 0+\infty = +\infty & f_{3}(1,1), c(1,0) + f_{3}(0,1), c(0,1) + f_{3}(1,0), c(1,1) + f_{3}(0,0) \\ (0,0) & (1,0) & (0,1) \\ f_{2}(1,0) = \min \begin{cases} 0+\infty = +\infty & f_{3}(0,0), c(1,0) + f_{3}(0,0) \\ (0,0) & (1,0) & (0,1) \\ f_{3}(0,2) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,2), c(0,1) + f_{3}(0,1) \\ (0,0) & (0,1) \\ f_{3}(0,1) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,2), c(0,1) + f_{3}(0,1) \\ (0,0) & (0,1) \\ f_{4}(0,1) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,2), c(0,1) + f_{3}(0,1) \\ (0,0) & (0,1) \\ f_{4}(0,1) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,2), c(0,1) + f_{3}(0,1) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = min \begin{cases} 0+\infty = +\infty & f_{4}\infty = +\infty \\ c(0,0) + f_{3}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = f_{4}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = f_{4}(0,0) \\ (0,0) & (0,1) \\ f_{4}(0,0) = f_{4}(0,0) \\ f_{4}(0,0) = f_{4}(0,0) \\ f_{4}(0,0) = f_{4}(0,0) \\ f_{4}(0,0) & f_{4}(0,0) \\ f_{4}(0,0) \\ f_{4}(0,0) & f_{4}(0,0) \\ f$$

Stage 1:

$$f_{1}(1,2) = \min \left\{ \begin{array}{c} 0+\infty = +\infty & 5+\infty = +\infty & 7+14 = 21 \\ c(0,0) + f_{2}(1,2), & c(1,0) + f_{2}(0,2), & c(0,1) + f_{2}(1,1), & c(1,1) + f_{2}(0,1) \\ (0,0) & (1,0) & (0,1) \end{array} \right\} = 2$$

=) Optimal solution: $x_1 = (1,1)$ $x_2 = (0,1)$ Build at location 1: 1000 oil capacity Build at location 2: 1000 gas capacity

Optimal value: 21 minimum total cost = \$21M

Note: $x_1 = (0, 1)$, $x_2 = (1, 1)$ is also an optimal solution